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(54) Title: GLOBAL SIGNAL DISTRIBUTION ARCHITECTURE IN A FIELD PROGRAMMABLE GATE ARRAY (57) Abstract A global signal distribution architecture for an FPGA architecture that has a plurality of multiplexers (80) with inputs and an output. The outputs of the plurality of multiplexers (80) are coupled to global I/O lines (16) that are coupled to a global signal distribution bus (18) spanning a highest level in the FPGA architecture. A plurality of switch matrices (106) are coupled to global signal distribution bus (18) and to a plurality of utility conductors (108) that are coupled to at least one multiplexer associated with a lowest level in the FPGA architecture.		

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SPECIFICATION

GLOBAL SIGNAL DISTRIBUTION ARCHITECTURE IN A FIELD PROGRAMMABLE GATE ARRAY

BACKGROUND OF THE INVENTION

1. Field Of The Invention

The present invention relates to global signal distribution architectures. More particularly, the present invention relates to global signal distribution architecture in a field programmable gate array (FPGA) that is reprogrammable.

2. The Prior Art

In the acceptance of an FPGA with capacity of up to a million gates, it is crucial that high speed fan-out nets be implemented in the FPGA design. High fanout nets in an FPGA are typically separated into four categories: global utilities, local clock/set/reset, control signals and high fanout data.

Examples of global utilities are clock, set or reset signals that define the main clock domains in the device by reaching many of the flip-flops in the FPGA. The local clock/set/reset signals, though they may have a medium to high fanout typically occur infrequently.

Well known examples of control signals include flip-flop enable and multiplexer select, however, control function can be more generally described as being orthogonal

to the data flow in the design. Such signals have a medium to high fanout and may occur frequently in a design. Another important characteristic of a control signal, is that the source of the control signal may originate in a different logic component such that the control signal driver may not be situated in the same physical hierarchy as the load of the control signal driver. The category of high fanout data are those high fanout signals that do not qualify as control signals.

In the design of high speed fanout nets the requirements of each of these four categories of high fanout signals must be taken into consideration. Accordingly, it is an object of the present invention to meet the requirement of all four categories of high fanout signals, while maintaining flexibility in the FPGA.

BRIEF DESCRIPTION OF THE INVENTION

According to the present invention, global utility signals that have their origin from input pins to the FPGA or from internal signals in the FPGA are distributed by a global distribution architecture to entities in the FPGA hierarchy known as clusters. The global distribution architecture includes I/O blocks that have I/O pins, buffers, boundary scan registers, interconnect to delay locked loops, and a global I/O (GIO) routing channel that includes global interconnect conductors which are coupled to the outputs of global multiplexers.

A global signal may be provided through a global multiplexer on the GIO routing channel from any of four locations, namely, an input buffer, the output of a BSR, from the output of a DLL, and from an internal FPGA signal. The GIO routing channels from the I/O blocks at both the top and bottom of the FPGA architecture are coupled into a

B16x16 tile of the FPGA architecture to form a global signal bus (GSB) in each of the B16x16 tiles of the FPGA architecture.

G3 routing channels run adjacent the GSB. The G3 routing channels can be coupled to selected interconnect conductors in the expressway routing channels M3.

As the G3 routing channels and the GSB traverse a B16x16 tile, they pass through G3MAT switch matrices that provide access to B4x4 tiles in the FPGA architecture. The G3 routing channels and the GSB are also coupled to user SRAM modules associated with a B16x16 tile.

A G3MAT switching matrix switches the signals on the GSB and a G3 routing channel onto a B4x4 utility routing channel. The GSB has GENERAL PURPOSE/ENABLE interconnect conductors, SET interconnect conductors, and CLOCK interconnect conductors.

At the lowest level in the FPGA architecture, the utilities from a G3MAT switch matrix are coupled to the inputs of four 8-input multiplexers labelled G, E, C, and S that select from the GENERAL PURPOSE signals, the ENABLE signals, the CLOCK signals and the SET signals. The outputs of multiplexers G, E, C, and S are coupled into each of the four clusters in the lowest level of the FPGA architecture on signal lines G, C, S, and E that correspond to the multiplexers labelled G, E, C, and S.

BRIEF DESCRIPTION OF THE DRAWINGS

FIG. 1 is a block diagram of a semi-hierarchical FPGA architecture with the global

signal distribution architecture according to the present invention

FIG. 2 is a block diagram of a B16x16 tile in an FPGA and the associated routing resources in the middle level of semi-hierarchical architecture according to the present invention.

FIG. 3 is a block diagram of a B2x2 tile in an FPGA and the connection of the routing resources in the lowest level to the middle level of a semi-hierarchical architecture according to the present invention.

FIG. 4 is a block diagram of a B2x2 tile in an FPGA and the routing resources in the lowest level of a semi-hierarchical architecture according to the present invention.

FIG. 5 is a schematic diagram of the I/O block depicted in FIG. 1 according to the present invention.

FIG. 6 illustrates the distribution of global signals to various circuits in an B16x16 tile by a global distribution bus according to the present invention.

FIG. 7 illustrates a switching matrix for a global signal bus and a G3 routing channel according to the present invention.

FIG. 8 illustrates the connection of a global signal bus and G3 routing channels to user SRAM blocks according to the present invention.

FIG. 9 illustrates the connection of utilities signals from a switching matrix to the clusters in a B1 block according to the present invention.

DETAILED DESCRIPTION OF A PREFERRED EMBODIMENT

Those of ordinary skill in the art will realize that the following description of the present invention is illustrative only and not in any way limiting. Other embodiments of the invention will readily suggest themselves to such skilled persons.

The present invention is implemented in a semi-hierarchical FPGA architecture having top, middle and low levels. According to the present invention, signals with high fanout may be distributed into the lowest level in the semi-hierarchical FPGA as global utility signals that have their origin from input pins to the FPGA at the highest level in the semi-hierarchical or from internal signals in the FPGA. To better understand the present invention, a description of the three levels of the semi-hierarchical architecture is made herein.

Turning now to FIG. 1 a block diagram of the top level of a semi-hierarchical architecture FPGA 10 with a global signal distribution architecture according to the present invention is illustrated. The top level of the architecture is an array of the B16x16 tiles 12 arranged in a rectangular array. A B16x16 tile 12 is a sixteen by sixteen array of B1 blocks. As will be described in detail below, a B16x16 tile 12 and its associated routing resources represents the middle level in the semi-hierarchical architecture, and a B1 block and its associated routing resources represents the lowest level in the semi-hierarchical architecture.

According to the present invention, the B16x16 tiles 12 are enclosed by I/O blocks 14 on the periphery. Each of the I/O blocks 14 include I/O pins, buffers, and boundary scan registers. Further, each of the I/O blocks 14 on the top and the bottom of the FPGA 10 include global I/O (GIO) routing channels 16-1 and 16-2, respectively, that preferably include sixteen interconnect conductors. The GIO routing channels 16 are coupled into each of the B16x16 tiles 12 to form a 32-bit global signal bus (GSB) 18 in each of the B16x16 tiles 12. The I/O blocks 14 on the top and the bottom of the FPGA 10, a GIO routing channel 16, and a GSB 18 will be described in greater detail below.

On each of the four sides of a B16x16 tile 12, and also associated with each of the I/O blocks 14 is freeway routing channel 20. It should be appreciated that on each side of a B16x16 tile 12 there are two freeway routing channels 20, either as a result of the disposition of two freeway routing channels 20 between adjacent B16x16 tiles 12 or as a result of the disposition of two freeway routing channels 20 between a B16x16 tile 12 and an adjacent I/O block 14.

It should be appreciated that the number of B16x16 tiles 12 in the FPGA 10 may be fewer or greater than the four shown in FIG. 1. The width of a freeway routing channel 20 in the FPGA 10 can be changed to accommodate different numbers of B16x16 tiles 12 without disturbing the internal structure of the B16x16 tiles 12. In this manner, the floorplan of the FPGA 10 can readily be custom sized by including the desired number of B16x16 tiles 12 in the design.

The freeway routing channels 20 can be extended in any combination of

directions at each end by a freeway turn matrix (F-turn) 22. An F-turn 22 is an active device that includes tri-state buffers and a matrix of reprogrammable switches. The reprogrammable switches are preferably SRAM pass devices. The interconnect conductors in the freeway routing channels 20 that are fed into an F-turn 22 may be coupled to many of the other interconnect conductors in the freeway routing channels 20 that come into the F-turn 22 by the programmable switches. Further, the interconnect conductors in the freeway routing channels 16 that are fed into an F-turn 22 continue in the same direction through the F-turn 22, even though the interconnect conductors are coupled to other interconnect conductors by the reprogrammable switches. A description of the implementation of an F-turn 22 is beyond the scope of this disclosure and will not be made herein to avoid overcomplicating the disclosure and thereby obscuring the present invention.

The freeway routing channels 20 along with the F-turns 22 form a course mesh. A freeway routing channel 20 will very rarely be utilized all by itself without any extension, since such distances are abundantly covered by the routing resources in the middle hierarchy to be described below. A freeway routing channel 20 is primarily intended to be used in conjunction with one or more other freeway routing channel 20 in any direction that together can span a distances of two or more B16x16 tiles 12.

The freeway routing channels 20 are hardwired to selected ones of the interconnect conductors in the GIO routing channels 16 by global-freeway (G-F) turns 24. This provides an additional path for routing signals between the freeway routing channels 20 and the GIO routing channels 16. The selection of coupling a freeway routing channel 20 to one of the hardwired connections in the G-F turn 24 is made by

the F-turn 22 that is adjacent the G-F turn 24.

In FIG. 2, a block diagram of a B16x16 tile 12 and the associated routing resources in the middle level of hierarchy is illustrated. The B16x16 tile 12 is a sixteen by sixteen array of B1 blocks 30. To avoid overcomplicating the drawing figure, only the B1 blocks 30 in a single row and a single column are indicated by the reference numeral 30. The B16x16 tile 12 is based on the repetition and nesting of smaller groupings (tiles) of B1 blocks 30. The smallest tile that is directly replicated and stepped is a B2x2 tile 32 that includes a two by two array of four B1 blocks 30. The B2x2 tiles 32 are stepped into a four by four array of sixteen B1 blocks 30 in a B4x4 tile 34, and the B4x4 tiles 34 are stepped into a eight by eight array of sixty-four B1 blocks 30 in a B8x8 tile 36. A B16x16 tile 12 includes four B8x8 tiles 36.

Though not depicted in FIG. 2, the B16x16 tile 12 further includes a block of user assignable static random access memory (SRAM), disposed between the two upper B8x8 tiles 36, and a block of SRAM disposed between the two lower B8x8 tiles 36. According to the present invention, the SRAM blocks will be described in greater detail below.

The routing resources in the middle level of hierarchy are termed expressway routing channels. There are three types of expressway routing channels, namely M1, M2, and M3. In FIG. 2, only a single row and a single column of expressway routing channels M1, M2, and M3 are denominated to avoid overcomplicating the drawing figure. In a preferred embodiment of the present invention, there is a single group of nine interconnect conductors in an M1 expressway routing channel, two groups of nine

interconnect conductors in an M2 expressway routing channel, and six groups of nine interconnect conductors in an M3 expressway routing channel.

The expressway routing channels M1, M2, and M3 are segmented so that each expressway routing channel M1, M2, and M3 spans a distance of a B2x2 tile 32, a B4x4 tile 34, and a B8x8 tile 36, respectively. Between each of the segments in the expressway routing channels M1, M2, and M3 are disposed extensions that can extend the expressway routing channel M1, M2, or M3 an identical distance along the same direction.

The extensions 38 that couple the segments in the expressway routing channels M1 and M2 are passive reprogrammable elements that are preferably an SRAM pass device. The extensions 38 provide a one-to-one coupling between the interconnect conductors of the expressway routing channels M1 and M2 on either side of the extensions 38. To avoid overcomplicating the drawing figure, only the extensions 38 in a single row and a single column are indicated by the reference numeral 38.

The segments of an M3 expressway routing channel are extended at the boundary of a B16x16 tile 12 where an expressway routing channel M3 crosses a freeway routing channel 20 by a freeway tab (F-tab) 40, and otherwise by an M3 extension 42. To avoid overcomplicating the drawing figure, only the F-tabs 40, and M3 extensions 42 in a single row and a single column are indicated by the reference numerals 40 and 42, respectively.

An F-tab 40 is an active device that includes tri-state buffers and a matrix of

reprogrammable switches. The reprogrammable switches are preferably an SRAM pass device. The interconnect conductors in the freeway routing channels 20 and the expressway routing channel M3 that are fed into an F-tab 40 may be coupled by the reprogrammable switches to many of the other interconnect conductors in the freeway routing channels 20 and the expressway routing channel M3 that come into the F-tab 40. A description of the implementation of the an F-tab 30 is beyond the scope of this disclosure and will not be made herein to avoid overcomplicating the disclosure and thereby obscuring the present invention.

Further, the interconnect conductors in the freeway routing channels 20 and the expressway routing channel M3 that are fed into an F-tab 40 continue in the same direction through the F-tab 40, even through the interconnect conductors are coupled to other interconnect conductors by the reprogrammable switches. When the F-tabs 40 are disposed between adjacent B16x16 tile 12, an expressway routing channel M3 continues on to another expressway routing channel M3. According to the present invention, as will be described in greater detail below, when the F-tabs 40 are disposed between a B16x16 tile 12 and an I/O block on either the top of the bottom of the FPGA 10, an expressway routing channel M3 is coupled into the I/O block 14.

Accordingly, an F-tab 40 implements the dual role of providing an extension of the middle level routing resources in a B16x16 tile 12 to the middle level routing resources in an adjacent B16x16 tile 12 or an I/O block 14, and providing access between the middle level routing resources of B16x16 tile 12 and a freeway routing channel 20 in the highest level of the architecture. An F-tab 40 can combine the two roles of access and extension simultaneously in the formation of a single net.

An M3 extension 42 is an active device that includes tristatable buffers coupled to a matrix of reprogrammable switches. The reprogrammable switches are preferably an SRAM pass device. The interconnect conductors in the expressway routing channel M3 that are fed into an M3 extension 42 may be coupled by the reprogrammable switches to many of the other interconnect conductors in the expressway routing channel M3 that come into the M3 extension 42. A description of the implementation of an M3 extension 42 is beyond the scope of this disclosure and will not be made herein to avoid overcomplicating the disclosure and thereby obscuring the present invention.

As depicted in FIG. 2, all of the expressway routing channels M1, M2, and M3 run both vertically through every column and horizontally through every row of B2x2 tiles 32. At the intersections of each of the expressway routing channels M1, M2, and M3 in the horizontal direction with the expressway routing channels M1, M2 and M3 in the vertical direction is disposed an expressway turn (E-turn) 44. To avoid overcomplicating the drawing figure, only the E-turns 44 disposed in the B2x2 tiles 22 in a single row and a single column are indicated by the reference numeral 44. A description of the implementation of an E-turn 44 is beyond the scope of this disclosure and will not be made herein to avoid overcomplicating the disclosure and thereby obscuring the present invention.

At the lowest level of the semi-hierarchical FPGA architecture, there are three types of routing resources, block connect (BC) routing channels, local mesh (LM) routing channels, and direct connect (DC) interconnect conductors. According to a

preferred embodiment of the present invention, there are nine interconnect conductors in each BC routing channel and six interconnect conductors in each LM routing channel. Of these three, the BC routing channels serve the dual purpose of being able to both couple B1 blocks 30 together at the lowest level in the architecture, and also provide access to the expressway-routing channels M1, M2, and M3 in the middle level of the architecture. In FIG. 3 aspects of the BC routing channels will be described, and in FIG. 4 aspects of the LM routing channels and the DC interconnect conductors will be described.

Turning now to FIG. 3, a B2x2 tile 32 including four B1 blocks 30 is illustrated. Associated with each of the B1 blocks 30 is a horizontal BC routing channel 50-1 and a vertical BC routing channel 50-2. Each horizontal BC routing channel 50-1 and vertical BC routing channel 50-2 is coupled to an expressway tab (E-tab) 52 to provide access for each B1 block 30 to the vertical and horizontal expressway routing channels M1, M2, and M3, respectively.

An E-tab 52 is an active device that includes tri-state buffers and a matrix of reprogrammable switches. The reprogrammable switches are preferably an SRAM pass device. The interconnect conductors in the BC routing channels 50 and the expressway routing channels M1, M2, and M3 that are fed into an E-tab 52 may be coupled by the programmable switches to many of the other interconnect conductors in the expressway routing channels M1, M2, and M3 that come into the E-tab 52. Further, the expressway routing channels M1, M2, and M3 that are fed into an E-tab 52 continue in the same direction through the E-tab 52, even through the interconnect conductors are coupled to other interconnect conductors by the reprogrammable

switches. A description of the implementation of an E-tab 52 is beyond the scope of this disclosure and will not be made herein to avoid overcomplicating the disclosure and thereby obscuring the present invention.

At the E-tabs 52, the signals provided on the BC routing channels 50 can connect to any of the expressway routing channels M1, M2, or M3. Once a signal emanating from a B1 block 30 has been placed on an expressway routing channel M1, M2 or M3 and traversed a selected distance, an E-tab 52 is employed to direct that signal onto a horizontal or vertical BC routing channel 50-1 or 50-2 into a B1 block 30 at a selected distance from the B1 block 20 from which the signal originated. As the connection between the routing resources at the lowest level in the architecture and the routing resources in the middle level of the architecture, the E-tabs 52 provide that the place and route of signals both inside and outside the B1 blocks 30 may be implemented independently from one another.

In FIG. 4, the expressway routing channels M1, M2, and M3 and the E-turn 44 have been omitted for clarity. As further depicted in FIG. 4, in addition to the horizontal and vertical BC routing channels 50-1 and 50-2 associated with each B1 block 30, there are also associated with each B1 block 20 four LM routing channels 54-1 through 54-4 and first and second DC interconnect conductors 56-1 and 56-2. The BC routing channels 50, the LM routing channels 54, and the DC interconnect conductors 56 provide significantly better performance than a strict hierarchy, and further help avoid congesting the expressway routing channels M1, M2, and M3. The BC routing channels 50 and the LM routing channels 54 combine to form two meshes. One is a mesh connection within a B1 block 20, and a second is a mesh connection between B1

blocks 20.

The BC routing channels 50 provide portions of the two meshes. In the portion of the mesh providing connection between adjacent B1 blocks 30, each horizontal and vertical BC routing channel 50-1 and 50-2 share an E-tab 52 with a horizontal or vertical BC routing channel 50-1 and 50-2 in an adjacent B1 block 30 that may be employed to couple a signal between adjacent B1 blocks 30 in a first direction. Further, each horizontal and vertical BC routing channel 50-1 and 50-2 share a BC extension 58 with a horizontal or vertical BC routing channel 50-1 and 50-2 in an adjacent B1 block 30 that may be employed to couple a signal between adjacent B1 blocks 20 in a second direction. The BC extensions 58 provide a one-to-one coupling between the interconnect conductors of the BC routing channels 50 on either side of the BC extensions 58. Accordingly, each BC routing channel 50, in the horizontal and vertical directions is coupled to the adjacent B1 blocks 30 in the corresponding horizontal and vertical directions by a E-tab 52 in a first direction along both the horizontal and vertical and in a second direction along both the horizontal and vertical by a BC extension 58.

From drawing FIG. 4, it should be appreciated that the LM routing channels 54-1 through 54-4 pass through the B1 block 30 as two vertical LM routing channels 54-1 and 54-4 and two horizontal LM routing channels 54-2 and 54-3, and that the intersections 60 of the vertical and horizontal LM routing channels 54 are hardwired along a diagonal.

The LM routing channels 54 also provide portions of the two meshes. In the portion of the mesh formed along with BC routing channels between B1 block 30, each

of the four LM routing channels 54-1 through 54-4 in each B1 block 30 shares an LM extension 62 with an LM routing channel 54-1 through 54-4 in an adjacent B1 block 30 in either the corresponding horizontal or vertical direction that may be employed to couple a signal between adjacent B1 blocks 30 in either the horizontal or vertical direction. The LM extensions 62 provide a one-to-one coupling between the interconnect conductors of the LM routing channels 54 on either side of the LM extensions 62. Accordingly, between adjacent B1 blocks 30 there are two LM routing channels 54 from each of the adjacent B1 blocks coupled by a LM extension 62 on all sides of adjacent B1 blocks 30.

The DC interconnect conductors 56-1 and 56-2 form a high performance direct connection between the logic elements in adjacent B1 blocks 30 to implement data path functions such as counters, comparators, adders and multipliers. As will be described below, each B1 block 20 includes four clusters of logic elements. Preferably, each of the four clusters includes two three input look-up tables (LUT3), a single two-input look-up table (LUT2), and a D-type flip-flop (DFF). In the DC interconnect conductor routing path, each of the DC interconnect conductors 56-1 and 56-2 is multiplexed to an input to a separate one of the two LUT3s in each of the four cluster of a B1 block 30. The DC interconnect conductors 56-1 and 56-2 are connected between vertically adjacent B1 blocks 30 as is illustrated in FIG. 4.

Turning now to FIG. 5, a portion of an I/O block 14, including a GIO routing channel 16 is depicted in greater detail. The portion of the I/O block 14 is shown above a B2x2 tile 32 that is disposed along the edge of a B16x16 tile 12. In FIG. 2, a B2x2 tile 32 disposed along the edge of a B16x16 tile 12 is illustrated with an expressway routing

channel M3 coupled to an F-tab 40. The interconnect conductors in the vertical BC routing channels 50 and the vertical LM routing channels 54 in the B2x2 tile 32 are also depicted. In FIG. 5, the F-tab 40 described with regard to FIG. 2, is distributed as the F-tabs 70-1 through 70-6, each of which is coupled to one of the six groups of nine interconnect conductors forming a single expressway routing channel M3. Further, as described above, between the B2x2 tile 32 on the edge of a B16x16 tile 12 and an I/O block 14 are disposed two freeway routing channels 20.

In the I/O block 14, input/output (I/O) pins 72 are coupled to input buffer 74 inputs and high impedance output buffer 76 outputs. The I/O pins 72 may carry clock, set, enable or global signals. However, the pins nearest the center of the FPGA typically have the lowest skew over the FPGA, and are therefore preferably reserved for the highest performance clock signals. The output of each of the input buffers 72 is coupled to the input of a boundary scan register (BSR) 78, and is also passed by the BSR 78 so that it is coupled to a first input of global multiplexer 80 associated with the GIO routing channel 16.

The output of each of the input buffers 74 may also be programmably coupled to a delay lock loop (DLL) reference line 80 or DLL feedback line 82 by reprogrammable elements 38 depicted by open circles. Only one of the reprogrammable elements is depicted by the reference numeral 84 to avoid overcomplicating the drawing figure. The reprogrammable elements 84, with one exception to be mentioned below, are preferably, SRAM pass devices, although those of ordinary skill in the art will readily appreciate that other types of programmable elements may also be suitably employed according to the present invention.

The input of each of the high impedance output buffers 76 may be programmably coupled to the DLL feedback line 82 by a reprogrammable element 84. The DLL feedback line 82 may also be programmably coupled to interconnect conductors in the BC routing channel 50 by programmable elements that are preferably tri-state buffers 86. The input of each of the high impedance output buffers 76 is also connected to one of the distributed F-tabs 70-1 through 70-5, and may also be coupled by a reprogrammable element to interconnect conductors in the LM routing channels 54 of the B2x2 tile 32. The F-tab 70-6 may be programmably coupled to either the output of each BSR 78 or to the enable input of each high impedance output buffer 76.

The output of each BSR 78 is also coupled to a second input of a separate global multiplexer 80 and to a separate F-tab 70-1 through 70-5. The DLL clock output line 88 is coupled to a third input of each of the global multiplexers 32, and each of the F-tabs 70-1 through 70-5 is coupled to a fourth input of a separate global multiplexer 80.

The outputs of the global multiplexers 80 are coupled to tristatable buffers 90 whose outputs are coupled to interconnect conductors in the GIO routing channels 16. Accordingly, from the above discussion it should be appreciated that at the output of the tristatable buffers 90 a global signal may be provided to the FPGA architecture 10 from any of four locations, namely, an input buffer 74, the output of a BSR 78, from the output of a DLL on the DLL output line 88, and from a signal in a B16x16 tile 12 via an F-tab 70-1 through 70-4.

It should be observed that more than one tristatable buffer 90 may drive a single

interconnect conductor in the GIO routing channel 16. The number of tristatable buffers 90 driving a single interconnect conductor in the GIO routing channel 16 depends on the number of B16x16 tiles 12 in the FPGA architecture 10. When the FPGA architecture 10 is 1, 2, 3, or 4 B16x16 tiles 12 wide, then the number of tristatable buffers 90 preferably coupled to a single interconnect conductor in the GIO routing channel 16 are 2, 4, 6, or 8, respectively. As depicted in FIG. 1, each of the GIO routing channels 16 from both the top and bottom of the FPGA architecture 10 are fed into a 32-bit GSB 32 in each of the B16x16 tiles 12.

Turning now to FIG. 6, the distribution of the global signals to one of four B8x8 tiles 36 in a B16x16 tile 12 by the GSB 18 is illustrated. It should be appreciated that the GSB 18 illustrated in FIG. 6 provides signals to each of the four B8x8 tiles 36 in a B16x16 tile 12 in a manner similar to the single B8x8 tile 36 that is illustrated. The B8x8 tile 36 includes four B4x4 tiles 34.

Disposed between the GSB 18 and each of the B8x8 tiles is a G3 routing channel 100. Preferably each G3 routing channel 100 has twelve interconnect conductors. The G3 routing channels 100 form intersections 102 with the expressway routing channels M3. Disposed at selected ones of the intersections 102 are reprogrammable elements that couple can couple interconnect conductors in G3 routing channels 100 with the interconnect conductors in the expressway routing channels M3 at the selected ones of the intersections 102. The reprogrammable elements are preferably SRAM pass devices.

At the freeway routing channels 20, selected interconnect conductors in the G3 routing channels 100 are coupled to global tabs (G-tabs) 104. A G-tab 104 is an active

device that includes tri-state buffers and a matrix of reprogrammable switches. The reprogrammable switches are preferably an SRAM pass device. The interconnect conductors in the freeway routing channels 20 and the G3 routing channel 100 that are fed into an G-tab 104 may be coupled by the reprogrammable switches to many of the other interconnect conductors in the freeway routing channels 20 and the G3 routing channel 100 that come into the G-tab 104. A description of the implementation of the a G-tab 104 is beyond the scope of this disclosure and will not be made herein to avoid overcomplicating the disclosure and thereby obscuring the present invention. Further, the interconnect conductors in the freeway routing channels 20 and the G3 routing channel 104 that are fed into a G-tab 104 continue in the same direction through the G-tab 104, even through the interconnect conductors are coupled to other interconnect conductors by the reprogrammable switches.

For each B4x4 tile 34, signals on a G3 routing channel 100 and the GSB 18 are coupled into the B4x4 tile 34 by a G3 switch matrix (G3MAT) 106 on a B4x4 utility routing channel 108. In FIG. 7, a G3MAT 106 is depicted in greater detail. The GSB 18 is preferably partitioned into sixteen GENERAL PURPOSE/ENABLE interconnect conductors 110, eight SET interconnect conductors 112, and eight CLOCK interconnect conductors 114. A G3 routing channel 100 is also illustrated. In the G3 MAT 106, the interconnect conductors in the G3 routing channel and the GSB 18 form intersections with the B4x4 utility routing channel 108. The interconnect conductors in the B4x4 utility routing channel 108 are buffered by tri-state buffers 116. Disposed at selected intersections are reprogrammable elements 118 depicted as open circles. To avoid overcomplicating the drawing figure only a single tri-state buffer 116, and a single reprogrammable element is indicated by the reference numeral 118.

The G3 routing channels 100 and GSB 18 are also coupled to user SRAM modules 120 by the SRAM switching block 122. As illustrated in FIG. 8, in the SRAM switching block 122, the eight CLOCK interconnect conductors 112 of the GSB 18 can be coupled to the write and read clock (WRCK, RCK) inputs of the user SRAM blocks 120, the sixteen GENERAL PURPOSE/ENABLE interconnect conductors 108 can be coupled to the write-enable and read-enable (WEN, REN) inputs of the user SRAM blocks 120, and the G3 routing channels 100 can be coupled to the WRCK/RCK inputs, and the WEN/REN inputs of the user SRAM blocks 120.

FIG. 9 illustrates a B4x4 utility routing channel 108 from a G3MAT 106 coupled into a B1 block 30 according to the present invention. The connection of the B4x4 utility routing channel 108 to the remaining B1 blocks 30 in the B4x4 tile 34 are similar to those depicted.

Each B1 block 30 includes four clusters 130-1 through 134-4 of devices. Within the B1 block 30, the BC routing channels 50 and the LM routing channels 54 have been omitted for to avoid overcomplicating the drawing figure. Each of the four clusters 130-1 through 130-4 includes first and second LUT3s 132-1 and 132-2, respectively, a LUT2 134, and a DFF 136. Each of the LUT3s 132 have first, second, and third inputs indicated as "A", "B", and "C", and a single output indicated as "Y". Each of the LUT2s 134 have first and second inputs indicated as "A" and "B", and a single output indicated as "Y". With a LUT3 132, any three input boolean logic function may be implemented, and with a LUT2 134 any two input boolean logic function may be implemented.

Each DFF 136 has a data input indicated as "D", an enable input indicated as "E", a clock input, a preset input, and a data output indicated as "Q". The DFF 136 may also be configured as a latch. In each of the clusters 130-1 through 130-4, the outputs "Y" of the LUT3s 132-1 and 132-2 are multiplexed to the input of DFF 136, and further multiplexed with the output "Q" of the DFF 136 to form first and second outputs of each of the clusters 130-1 through 130-4.

The B4x4 utility routing channel 108 from a G3MAT 106 is coupled to the inputs of four 8-input multiplexers 138, 140, 142, and 144 that select from the GENERAL PURPOSE/ENABLE signals 110, the CLOCK signals 114 and the SET signals 112, respectively. As such, the multiplexers 138, 140, 142, and 144 are labelled G, E, C, and S. It should be noted that for the remaining B1 blocks 30 in the B2x2 tile 32, the multiplexers 142 and 144 may be employed, but multiplexers similar to the multiplexers 138 and 140 must be provided for each of the remaining B1 blocks 30 in the B2x2 tile 32.

The outputs of multiplexers 138, 140, 142, and 144 are coupled to tri-state buffers 146, and the outputs of the tri-state buffers 146 are coupled into each of the four clusters 130-1 through 130-4 on the signal lines 148, conveniently depicted as G, C, S, and E to correspond to the multiplexers 138, 140, 142, and 144 to which they are coupled.

In each of the clusters 130-1 through 130-4, the output of the multiplexer 138 on signal line G is coupled to the A input of LUT3 132-1 and to the B input of LUT3 132-2. The output of the multiplexer 140 on signal line E is coupled to the B input of LUT3

132-1 and to the A input of LUT3 132-2. The output of the multiplexer 142 on signal line C is coupled to a multiplexer whose output is coupled to the C input of the LUT3 132-1 in clusters 130-1 and 130-2 and the C input of the LUT3 132-2 in clusters 130-3 and 130-4. The output of the multiplexer 144 on signal line S is coupled to a multiplexer whose output is coupled to the S input of the LUT3 132-1 in clusters 130-3 and 130-4 and the C input of the LUT3 132-2 in clusters 130-1 and 130-2. Finally, the outputs of multiplexers 140, 142, and 144 (multiplexers E, C, and S) are coupled to the enable, preset, and clock inputs of the DFF 136 in each of the clusters.

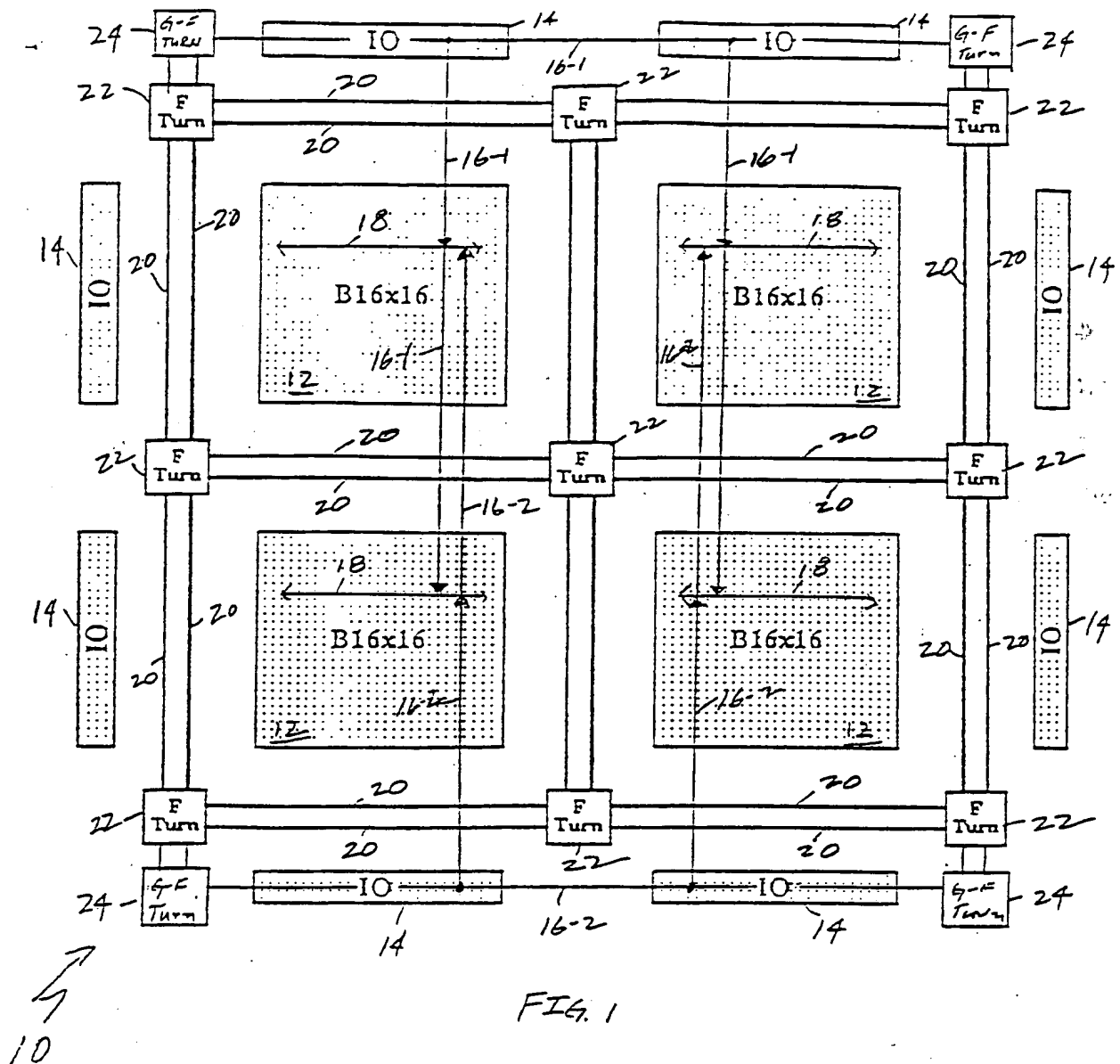
Each of the DC interconnect conductors 56-1 and 56-2 is multiplexed in a serial fashion with the C and S lines to the "C" inputs of LUT3s in each cluster 130-1 through 130-4. For example, in the serial connection, the DC interconnect conductor 56-1 is multiplexed with the C line to the "C" input of the LUT3 132-1 of the cluster 130-1. Next, the "Y" output of the LUT3 132-1 in cluster 130-1 is multiplexed with the S line to the "C" input of the LUT3 132-2 in cluster 130-2. Next, the "Y" output of the LUT3 132-2 in cluster 130-2 is multiplexed with the S line to the "C" input of the LUT3 132-1 in cluster 130-3. Next, the "Y" output of the LUT3 132-1 in cluster 130-3 is multiplexed with the C line to the "C" input of the LUT3 132-2 in cluster 130-4. Finally, the "Y" output of the LUT3 132- in cluster 130-4 passes out of the B1 block 20, and is multiplexed to the "C" input of the LUT3 132-1 in cluster 130-1 of the B1 block 20 disposed vertically below. The DC interconnect conductors 56-2 is similarly connected.

While embodiments and applications of this invention have been shown and described, it would be apparent to those skilled in the art that many more modifications

than mentioned above are possible without departing from the inventive concepts herein. The invention, therefore, is not to be restricted except in the spirit of the appended claims.

What is Claimed is:

1. — A global signal distribution architecture for an FPGA architecture comprising:
 - a plurality of multiplexers, each of said plurality of multiplexers having a plurality of inputs and an output;
 - a plurality of global I/O lines, each of said plurality of global I/O lines coupled to a separate output of one of said plurality of multiplexers;
 - a global signal distribution bus coupled to said plurality of global I/O lines; said global signal distribution bus spanning a highest level in the FPGA architecture;
 - a plurality of switch matrices, each of said switch matrices coupled to said global signal distribution bus to a plurality of utility conductors; and
 - at least one multiplexer associated with a lowest level in the FPGA architecture coupled to said plurality of utility conductors.



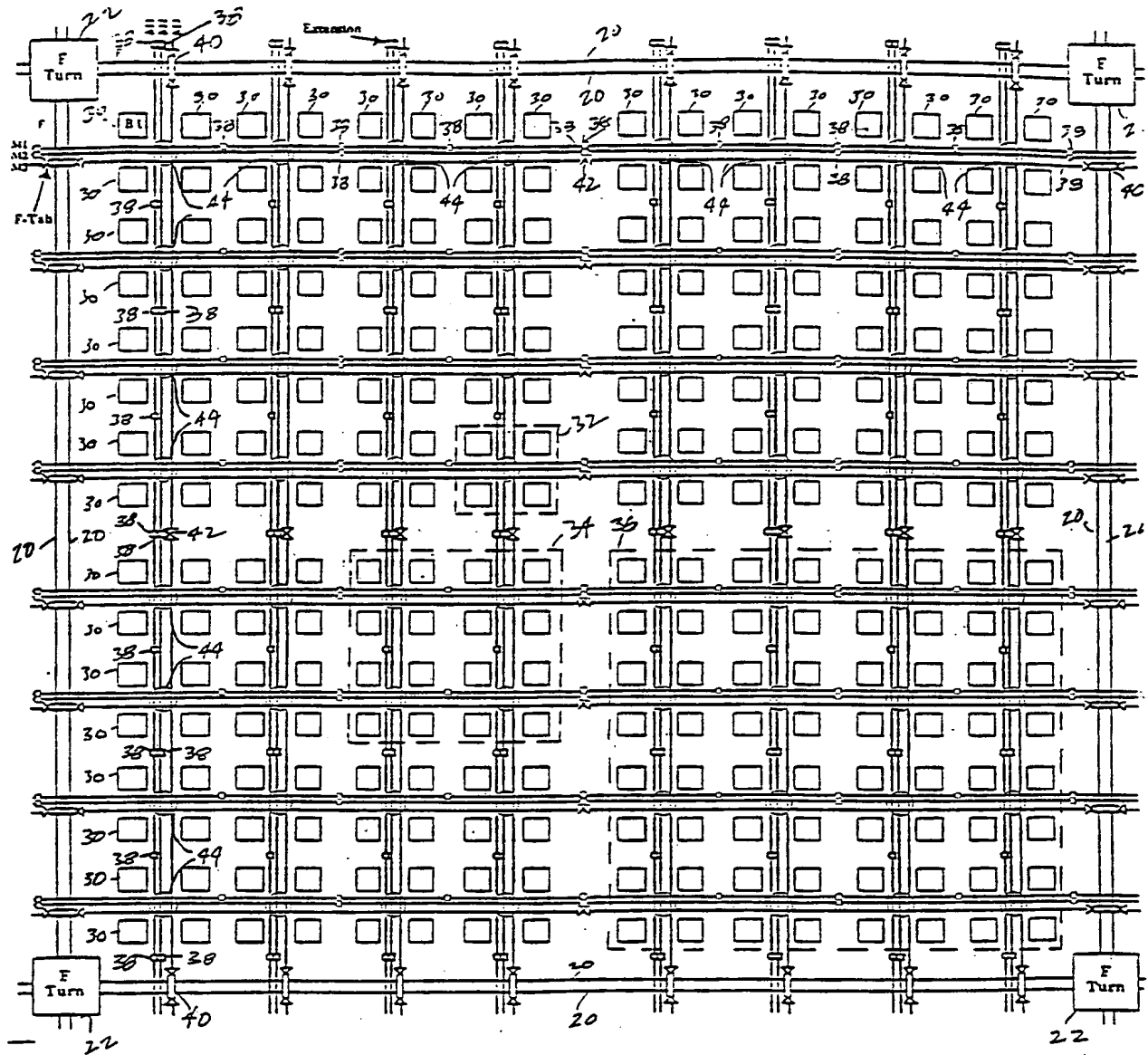


FIG. 2

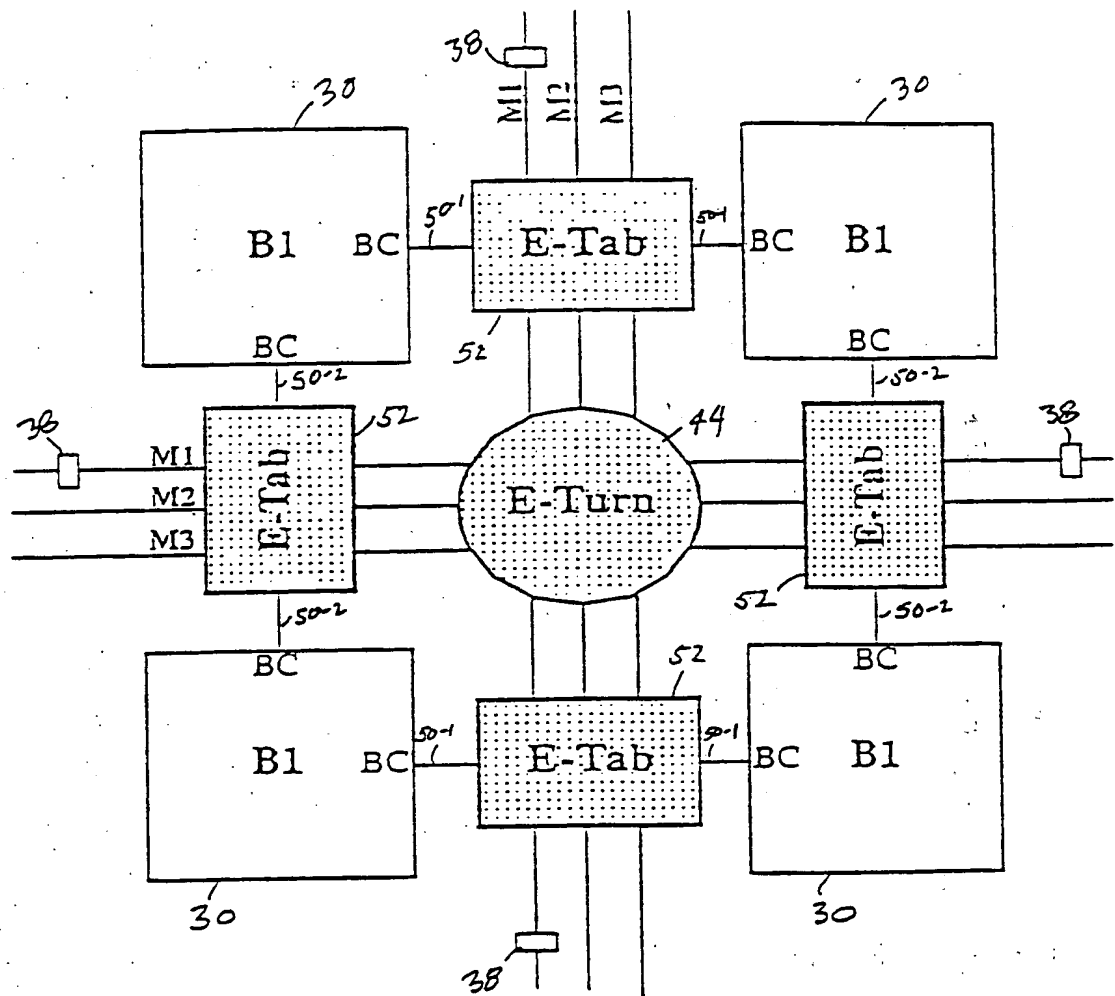


FIG. 3

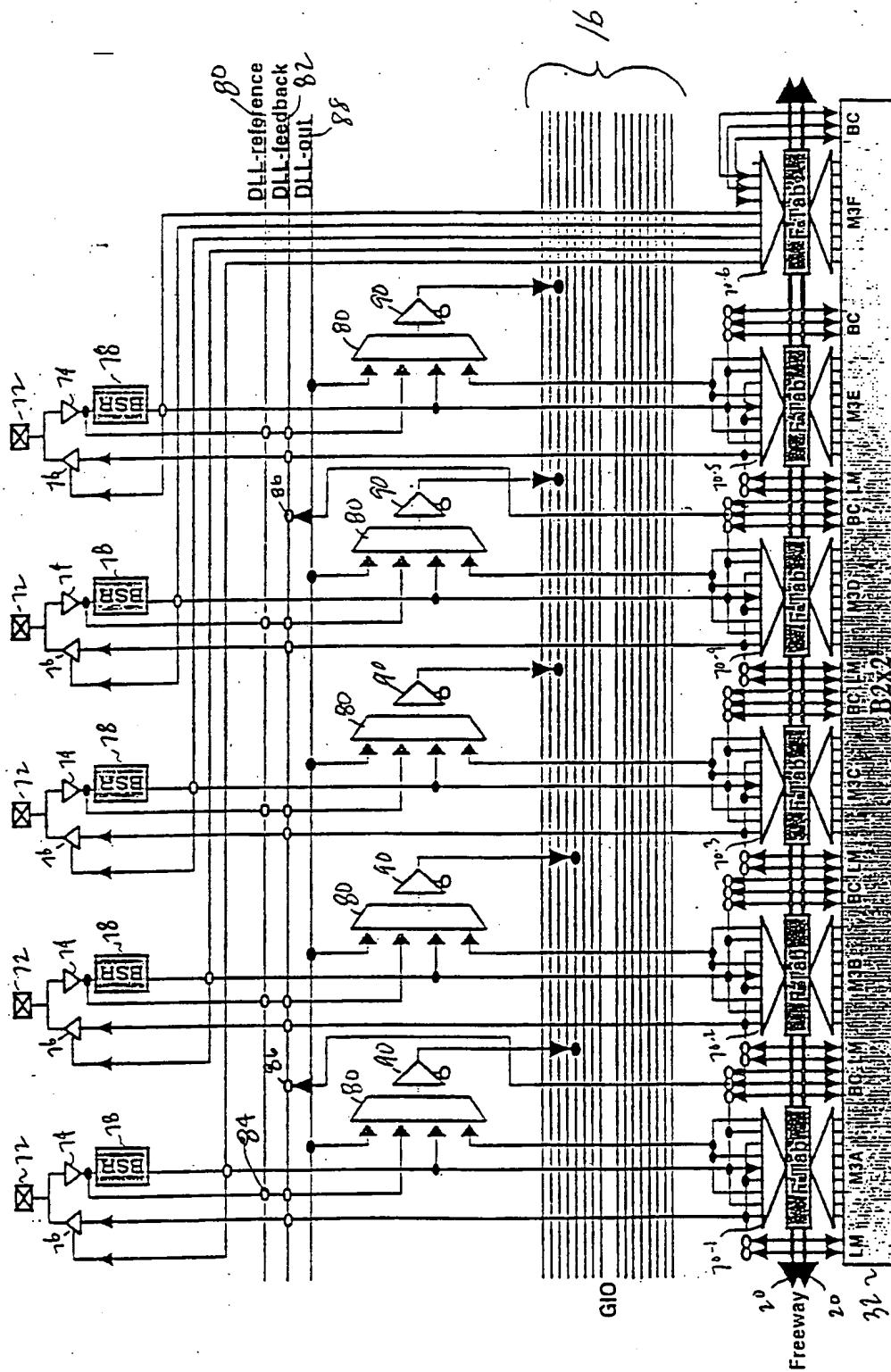


FIG. 5

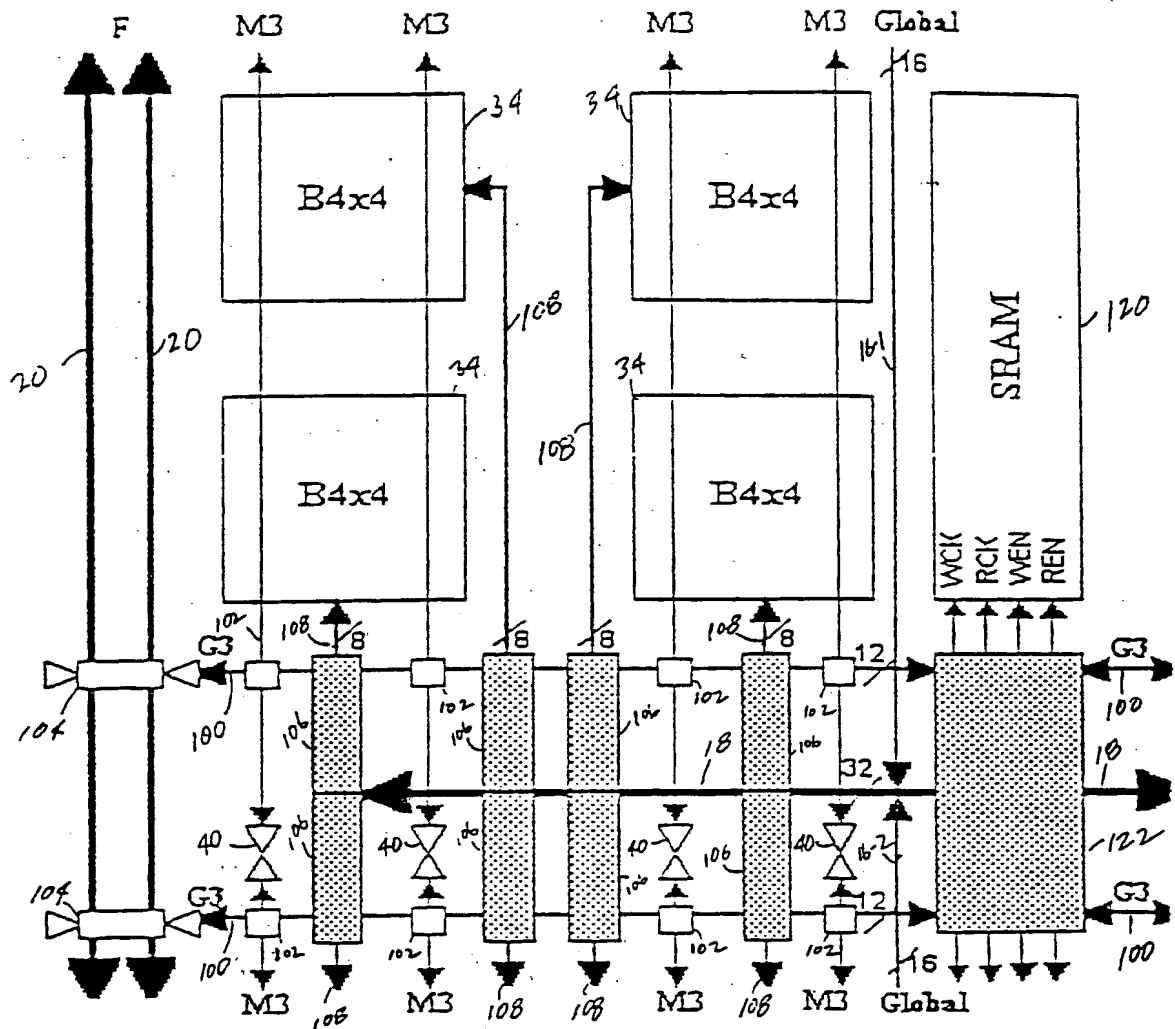
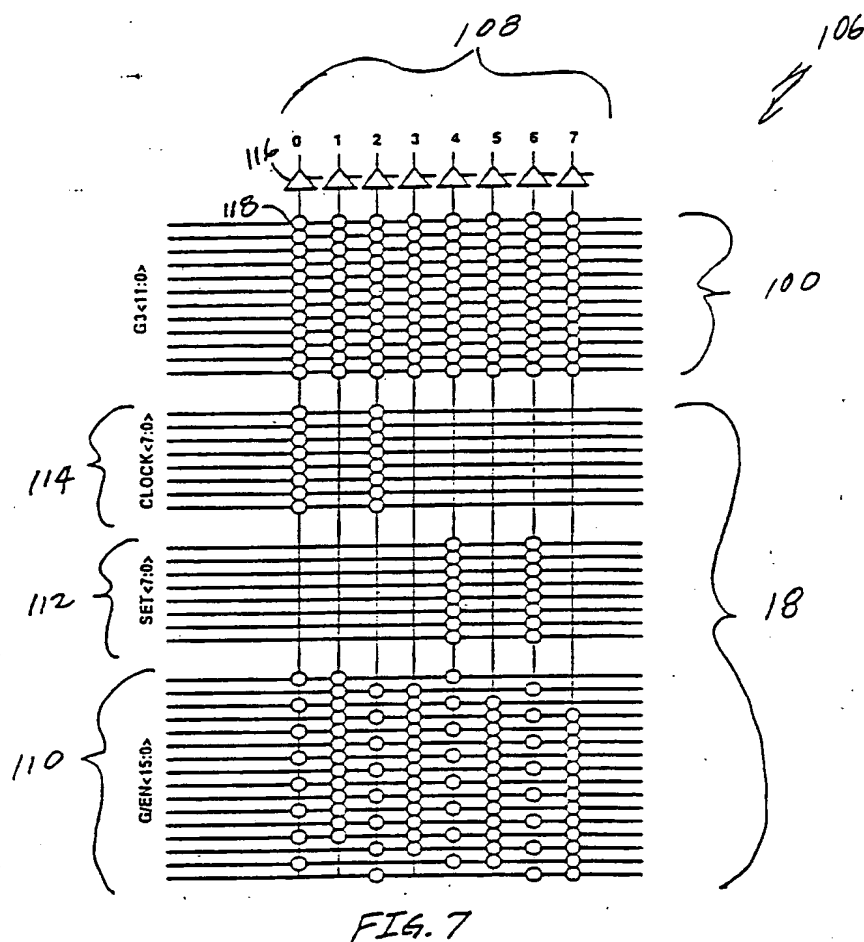


FIG. 6



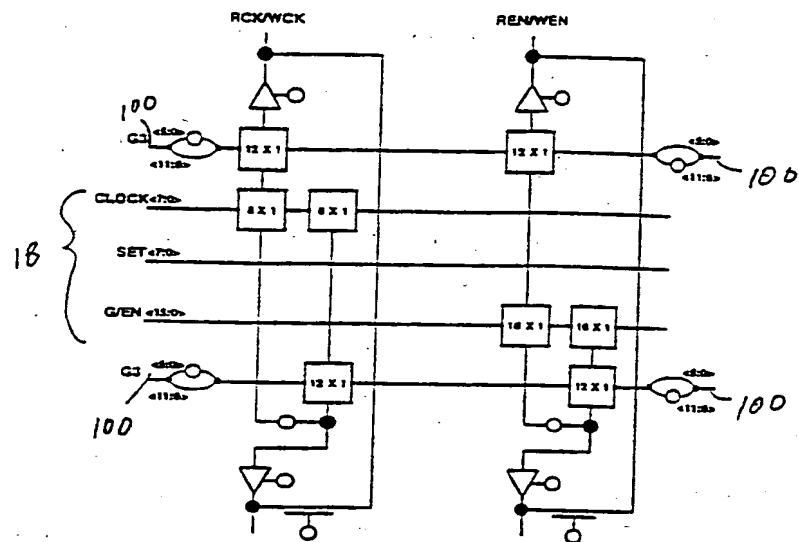


FIG. 8

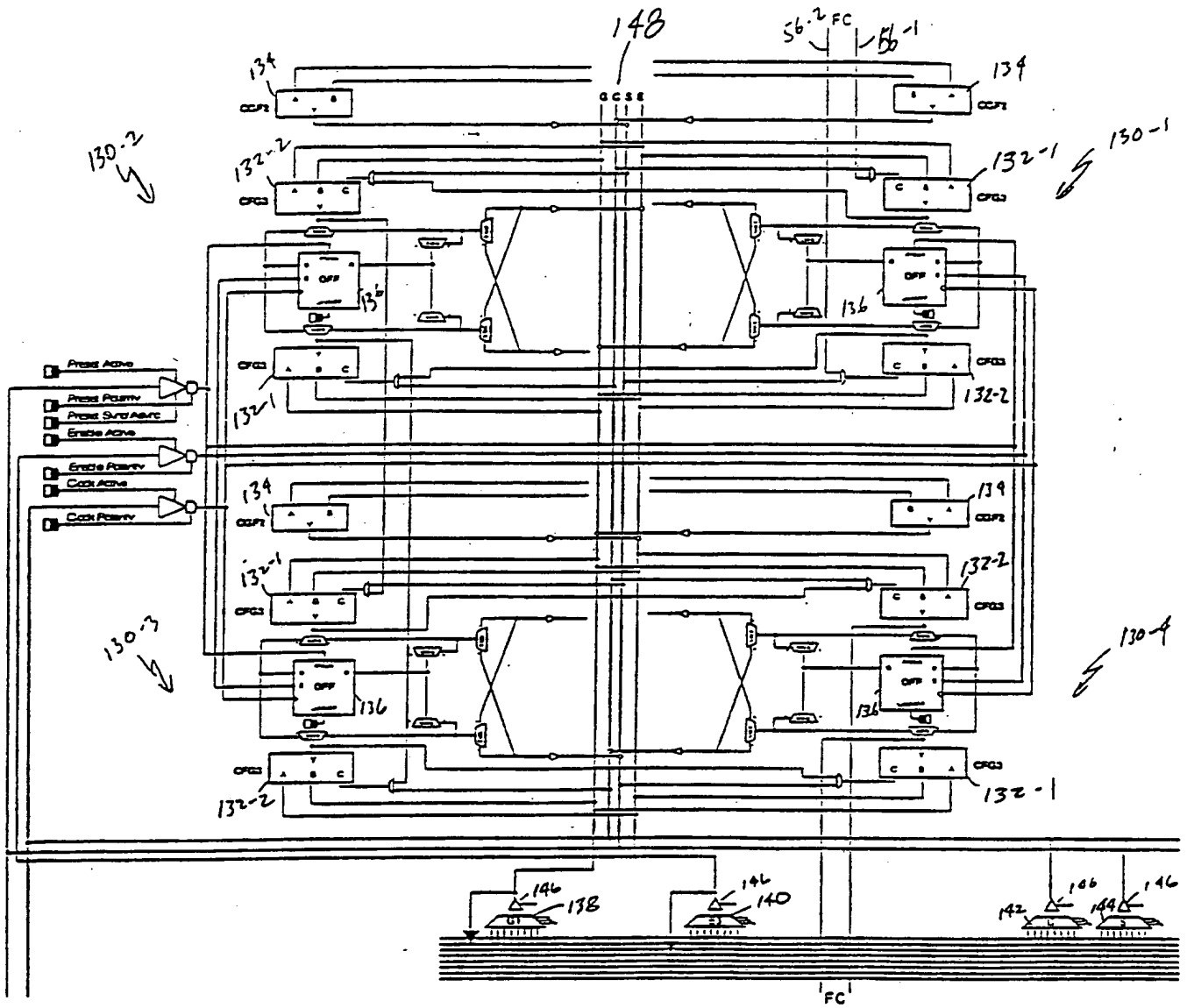


FIG. 9

INTERNATIONAL SEARCH REPORT

International application No.
PCT/US00/04477

A. CLASSIFICATION OF SUBJECT MATTER

IPC(7) :H03K 19/177; H01L 25/00

US CL :326/41, 40, 39, 38

According to International Patent Classification (IPC) or to both national classification and IPC

B. FIELDS SEARCHED

Minimum documentation searched (classification system followed by classification symbols)

U.S. : 326/41, 40, 39, 38, 37

Documentation searched other than minimum documentation to the extent that such documents are included in the fields searched
None

Electronic data base consulted during the international search (name of data base and, where practicable, search terms used)

USPTO APS BRS

Search terms: FPGA architecture, global routing for FPGA

C. DOCUMENTS CONSIDERED TO BE RELEVANT

Category*	Citation of document, with indication, where appropriate, of the relevant passages	Relevant to claim No.
A,P	US 5,942,914 A (REDDY et al.) 24 August 1999 (24-08-1999), abstract.	1
A	US 5,815,004 A (TRIMBERGER et al.) 29 September 1998 (29-09-1998), Figure 3.	1
X	US 5,455,525 A (HO et al.) 03 October 1995 (03-10-1995), Figures 1, 2, and 5.	1
A	US 5,598,109 A (LEONG et al.) 28 January 1997 (28-01-1997), Figures 2-5.	1

☐ Further documents are listed in the continuation of Box C.

☐ See patent family annex.

* Special categories of cited documents:

"A" document defining the general state of the art which is not considered to be of particular relevance

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later document published after the international filing date or priority date and not in conflict with the application but cited to understand the principle or theory underlying the invention

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document of particular relevance; the claimed invention cannot be considered novel or cannot be considered to involve an inventive step when the document is taken alone

"Y"

document of particular relevance; the claimed invention cannot be considered to involve an inventive step when the document is combined with one or more other such documents, such combination being obvious to a person skilled in the art

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document member of the same patent family

Date of the actual completion of the international search

02 MAY 2000

Date of mailing of the international search report

15 MAY 2000

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